

Everything You Wanted  
To Know  
About FIRST\_ROWS\_n  
But Were Afraid To Ask

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# Who am I

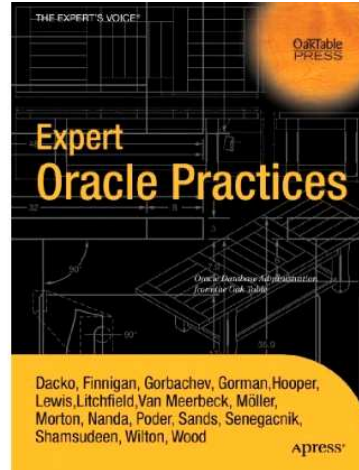
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- Independent Consultant
- Located in Germany
- Oracle ACE, OCP 8i, 9i, 10g
- Member of the OakTable Network
- My Blog: Oracle related stuff



# Who am I

Co-author of  
the forthcoming  
OakTable book  
"Expert  
Oracle  
Practices"



# Highlights

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- Introduction to the Cost-Based Optimizer's (CBO) OPTIMIZER\_MODEs
- Evolution of the CBO – FIRST\_ROWS, FIRST\_ROWS\_n
- Top N / Pagination queries, advanced techniques
- A deeper look at FIRST\_ROWS\_n, examples, quirks, oddities
- Questions and Answers

# Everything You Wanted To Know About FIRST\_ROWS\_n

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Introduction to the Cost-Based  
Optimizer's (CBO)  
OPTIMIZER\_MODEs

## Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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- CHOOSE between Rule Based Optimizer (RBO) and CBO: Default in pre-10g
- ALL\_ROWS: Default in 10g and later
- FIRST\_ROWS: Deprecated since 9i
- Oracle 9i introduced FIRST\_ROWS\_1, FIRST\_ROWS\_10, FIRST\_ROWS\_100 and FIRST\_ROWS\_1000 along with the FIRST\_ROWS(n) hint

## Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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When to use which mode?

- Often heard:

“In OLTP environments use  
FIRST\_ROWS / FIRST\_ROWS\_n,  
because typically you retrieve only a  
few rows with the queries”

# Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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Philosophy ALL\_ROWS:

The optimizer assumes that all rows from the generated result set will be processed / fetched by the client

This is true in most cases, therefore ALL\_ROWS is a good general approach and default from 10g on

# Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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## Philosophy ALL\_ROWS

- What about the “OLTP environment” statement?
- Assuming that the optimizer estimates the correct cardinality, queries identifying only a few rows still retrieve **all** of these few rows, so ALL\_ROWS is applicable here as well, no need to use FIRST\_ROWS/FIRST\_ROWS\_n

## Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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So when do we need FIRST\_ROWS(\_n)?

Philosophy FIRST\_ROWS(\_n):

- The optimizer assumes that only a ***few*** first rows from a ***larger*** result set will be processed / fetched by the client

## Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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In most cases these are so called Top N /  
Pagination queries:

- Show the top N rows from an ordered result set. Popular examples: Search results in Google / Amazon
- Retrieve the result set in pages -  
Pagination queries
- This is only reasonable if a deterministic order has been defined for the result set

## Cost-Based Optimizer (CBO) - OPTIMIZER\_MODEs

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- If you don't retrieve only the first rows from a larger result set but the `FIRST_ROWS(_n)` modes speed up the query processing, you very likely only were lucky due to the side effects of the `FIRST_ROWS(_n)` modes
- You should be able to get the same performance with `ALL_ROWS` provided that the optimizer estimates are in the right ballpark

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Later we will see that the optimizer will still get it right when the query returns only a few rows, all of them are fetched/processed, the estimates are in the right ballpark and the `FIRST_ROWS(n)` mode gets used with  $n \geq$  number of rows returned by the query

# Everything You Wanted To Know About FIRST\_ROWS\_n

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Evolution of the CBO –  
FIRST\_ROWS,  
FIRST\_ROWS\_n

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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- FIRST\_ROWS introduced with the CBO in Version 7
- FIRST\_ROWS uses a mixture of heuristics and cost based optimization
- Some built-in high-level constraints are:
  - Use index to avoid SORT operation
  - Try to avoid Hash / Sort Merge Join, use NESTED LOOP join

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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## Example Setup (1)

```
create table first_k_rows_test
(
  sort_order  varchar2(10) null,
  filter_col  number       null,
  filler      char(200)
);

insert into first_k_rows_test
select dbms_random.string('x', 10) as sort_order,
       case when level <= 10 then 1 else 0 end as filter_col,
       'x' as filler
from dual
connect by level <= 100000;

commit;
```

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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## Example Setup (2)

```
create index first_k_rows_test_idx1 on first_k_rows_test(filter_col);  
  
create index first_k_rows_test_idx2 on first_k_rows_test(sort_order);  
  
exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',  
method_opt=>'for all columns size 1 for columns filter_col size 254');
```

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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## ALL\_ROWS Baseline

```
select /*+ all_rows */
      *
from
      first_k_rows_test a,
      first_k_rows_test b
where
      a.filter_col = 1
and    a.sort_order = b.sort_order
order by
      a.sort_order;
```

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

## ALL\_ROWS Baseline Plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	23 (5)
1	SORT ORDER BY		10	23 (5)
2	NESTED LOOPS		10	22 (0)
3	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	10	2 (0)
* 4	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX1	10	1 (0)
5	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	1	2 (0)
* 6	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX2	1	1 (0)

Predicate Information (identified by operation id):

4 - access("A"."FILTER\_COL"=1)  
6 - access("A"."SORT\_ORDER"="B"."SORT\_ORDER")

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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## FIRST\_ROWS query

```
select /*+ first_rows */
      *
from
      first_k_rows_test a,
      first_k_rows_test b
where
      a.filter_col = 1
and    a.sort_order = b.sort_order
order by
      a.sort_order;
```

# Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

## FIRST\_ROWS plan

```
-----  
| Id | Operation                               | Name                               | Rows | Cost (%CPU)|  
-----  
| 0 | SELECT STATEMENT                         |                                     | 1    | 100K (1)|  
| 1 | NESTED LOOPS                             |                                     | 1    | 100K (1)|  
| * 2 | TABLE ACCESS BY INDEX ROWID            | FIRST_K_ROWS_TEST                 | 1    | 100K (1)|  
| 3 | INDEX FULL SCAN                          | FIRST_K_ROWS_TEST_IDX2            | 100K | 432 (2)|  
| 4 | TABLE ACCESS BY INDEX ROWID            | FIRST_K_ROWS_TEST                 | 1    | 2 (0)|  
| * 5 | INDEX RANGE SCAN                        | FIRST_K_ROWS_TEST_IDX2            | 1    | 1 (0)|  
-----
```

Predicate Information (identified by operation id):

```
-----  
2 - filter("A"."FILTER_COL"=1)  
5 - access("A"."SORT_ORDER"="B"."SORT_ORDER")
```

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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- These high-level constraints can lead to inefficient execution plans:
  - Favoring an inefficient index over an highly selective index
  - Using inefficient NESTED LOOP joins where other join methods can be more efficient
- `_sort_elimination_cost_ratio` parameter influences FIRST\_ROWS calculation

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In previous releases (< 9.2.0.8) the FIRST\_ROWS mode by default was avoiding the sort at all costs - in recent releases this has changed, but you can still get rather inefficient plans albeit the change - the gory details follow in the last part

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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- FIRST\_ROWS\_n introduced in Oracle 9i, FIRST\_ROWS deprecated
- FIRST\_ROWS\_n: Fully cost based, no built-in heuristics (?)
- However: Still some internal tweaks that favor index access under certain circumstances, so it doesn't choose always the plan with the lowest cost

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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FIRST\_ROWS\_n approach:

- The optimizer knows about the requested number of rows (the “n” in FIRST\_ROWS\_n)
- Either explicitly set via OPTIMIZER\_MODEs:
  - FIRST\_ROWS\_[1 | 10 | 100 | 1000]
  - or via FIRST\_ROWS(n) hint: e.g. FIRST\_ROWS(50)

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## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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FIRST\_ROWS\_n approach:

- Or implicitly via ROWNUM predicates:  
SELECT  
FROM (query [ORDER BY order])  
WHERE ROWNUM <= n
- “\_optimizer\_rownum\_pred\_based\_fkr”  
parameter applies, default “true” in  
recent releases

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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FIRST\_ROWS\_n approach:

- Knowing the “n” is essential, because the optimizer now works out how many rows the result set is estimated to produce
- This is basically done performing an initial (and partial) ALL\_ROWS optimization

## Cost-Based Optimizer (CBO) - FIRST\_ROWS(\_n) Evolution

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FIRST\_ROWS\_n approach:

- With the knowledge of “n” and the total number of rows estimated, the optimizer can calculate a “proration factor”, which is basically

$$n / \text{total\_number\_of\_rows}$$

- This “proration factor” is then used in the following optimization steps when calculating costs

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# Everything You Wanted To Know About FIRST\_ROWS\_n

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Top N / Pagination queries,  
advanced techniques

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Modified example setup

```
truncate table first_k_rows_test;

insert into first_k_rows_test
  select dbms_random.string('x', 10),
         case when level <= 5000 then 1 else 0 end,
         'x' as filler
  from dual
  connect by level <= 100000;

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',
method_opt=>'for all columns size 1 for columns filter_col size 254');
```

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Simple Top N query

```
select
    *
from (
    select
        t.*
    from
        first_k_rows_test t
    order by
        sort_order
)
where
    rownum <= 10;
```

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We expect to get only 10 rows and we have an index in place on SORT\_ORDER.  
Let's examine the execution plan

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan - first attempt

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	5516 (1)
* 1	COUNT STOPKEY			
2	VIEW		100K	5516 (1)
* 3	SORT ORDER BY STOPKEY		100K	5516 (1)
4	TABLE ACCESS FULL	FIRST_K_ROWS_TEST	100K	833 (2)

Predicate Information (identified by operation id):

- 1 - filter(ROWNUM<=10)
- 3 - filter(ROWNUM<=10)

Index doesn't get used - why?

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Example Setup

```
create table first_k_rows_test
(
  sort_order  varchar2(10) null,
  filter_col  number      null,
  filler     char(200)
);
```

**B\*tree indexes don't cover expressions  
that consist entirely of NULLs!**

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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B\*tree index can only be used for ordering if the index covers all values (NULLs!)

- Either define column as NOT NULL or have predicates in place that imply NOT NULL
- Or make the index cover NULL values e.g. contains a NOT NULL column or use function-based index:

```
create index first_k_rows_test_idx2 on first_k_rows_test(  
sort_order,  
' ');
```

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan - second attempt

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	12 (0)
* 1	COUNT STOPKEY			
2	VIEW		10	12 (0)
3	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	100K	12 (0)
4	INDEX FULL SCAN	FIRST_K_ROWS_TEST_IDX2	10	2 (0)

Predicate Information (identified by operation id):

1 - filter(ROWNUM<=10)

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Modified example setup

```
drop index first_k_rows_test_idx1;  
  
drop index first_k_rows_test_idx2;  
  
create unique index first_k_rows_test_idx on first_k_rows_test  
(  
    filter_col  
, sort_order  
);
```

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Top N query including filter

```
select
    *
from (
    select
        t.*
    from
        first_k_rows_test t
    where
        filter_col = 1
    order by
        sort_order
)
where
    rownum <= 10;
```

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We expect to get only 10 rows and we have the ideal index in place supporting both filter and sort. Let's examine the execution plan

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan - first attempt

```
-----  
| Id | Operation                | Name                | Rows  | Cost (%CPU)|  
-----  
|  0 | SELECT STATEMENT         |                    |      10 | 1056 (2)|  
|*  1 |   COUNT STOPKEY         |                    |      |      |  
|  2 |     VIEW                 |                    |  4699 | 1056 (2)|  
|*  3 |       SORT ORDER BY STOPKEY|                    |  4699 | 1056 (2)|  
|  4 |         COUNT           |                    |      |      |  
|*  5 |           TABLE ACCESS FULL | FIRST_K_ROWS_TEST |  4699 |  834 (2)|  
-----
```

Predicate Information (identified by operation id):

```
-----  
1 - filter(ROWNUM<=10)  
3 - filter(ROWNUM<=10)  
5 - filter("FILTER_COL"=1)
```

Index again doesn't get used - why?

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Example Setup

```
create table first_k_rows_test
(
    sort_order    varchar2(10) null,
    filter_col    number          null,
    filler char(200)
);
```

Sorting character strings:

***Client*** NLS settings are relevant!

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# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## NLS sorting issues

- Set NLS\_SORT to BINARY in session to support sort by conventional index

```
ALTER SESSION SET NLS_SORT = BINARY;
```

- Create function-based index using NLSSORT function

```
create index first_k_rows_test_idx on first_k_rows (  
  filter_col  
  , NLSSORT(sort_order, 'NLS_SORT=GERMAN')  
);
```

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If you happen to have different NLS settings per client you theoretically need several NLSSORT based indexes - to support each different NLS based sort order

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan - second attempt

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	13 (0)
* 1	COUNT STOPKEY			
2	VIEW		11	13 (0)
3	COUNT			
4	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	11	13 (0)
* 5	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX		2 (0)

Predicate Information (identified by operation id):

1 - filter(ROWNUM<=10)  
5 - access("FILTER\_COL"=1)

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Pagination query using ROWNUM

```
select
  *
from (
  select
    rownum as rn
  , v1.*
  from (
    select
      t.*
    from
      first_k_rows_test t
    where
      filter_col = 1
    order by
      sort_order
  ) v1
  where
    rownum <= 20
  ) v2
where rn >= 11
order by rn;
```

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Extending the Top N query to a "Pagination" query - show the ordered result set in pages. The application requests a specific page.

It is obvious that using ROWNUM this is a bit cumbersome to write, since due to the nature how ROWNUM gets evaluated we need to write two levels of inline views.

Note that although you expect the result to be already correctly ordered if the index gets used, you still need to specify the final sort order to have the order of the result set guaranteed.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Pagination query using ROWNUM

```
select
  *
from (
  select
    rownum as rn
  , v1.*
  from (
    select
      t.*
    from
      first_k_rows_test t
    where
      filter_col = 1
    order by
      sort_order
    ) v1
  where
    rownum <= 20
  ) v2
where rn >= 11
order by rn;
```

This corresponds to the Top N query, but it needs now to identify as many rows as indicated by the page requested

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This part is the already shown Top N query

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Pagination query using ROWNUM

```
select
  *
from (
  select
    rownum as rn
  , v1.*
  from (
    select
      t.*
    from
      first_k_rows_test t
    where
      filter_col = 1
    order by
      sort_order
  ) v1
  where
    rownum <= 20
  ) v2
where rn >= 11
order by rn;
```

Now we discard all the rows except for those left for the page to display

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We filter all rows that don't belong to the requested page.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		20	23 (5)
1	SORT ORDER BY		20	23 (5)
* 2	VIEW		20	22 (0)
* 3	COUNT STOPKEY			
4	VIEW		20	22 (0)
5	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	20	22 (0)
* 6	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX	2	(0)

Predicate Information (identified by operation id):

- 2 - filter("RN">=11)
- 3 - filter(ROWNUM<=20)
- 6 - access("FILTER\_COL"=1)

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Note that the optimizer doesn't recognize the "pagination" - it estimates a final cardinality of 20.

Optionally you could wrap the pagination query into a third view requesting on the first "pagesize" rows via the ROWNUM predicate - this will make sure that the optimizer recognizes the page size requested.

The same will be achieved by using the FIRST\_ROWS(n) hint where n equals the pagesize.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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Same query, use FIRST\_ROWS(n) hint

```
select /*+ first_rows(10) */
      *
from   (
       select
         rownum as rn
       , v1.*
       from   (
              select
                t.*
              from
                first_k_rows_test t
              where
                filter_col = 1
              order by
                sort_order
            ) v1
       where
         rownum <= 20
      ) v2
where rn >= 11
order by rn;
```

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Here is the same query using the FIRST\_ROWS(10) hint - note the subtle difference in the execution plan

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	13 (8)
1	SORT ORDER BY		10	13 (8)
* 2	VIEW		10	12 (0)
* 3	COUNT STOPKEY			
4	VIEW		10	12 (0)
5	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	10	12 (0)
* 6	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX	2	(0)

Predicate Information (identified by operation id):

- 2 - filter("RN">=11)
- 3 - filter(ROWNUM<=20)
- 6 - access("FILTER\_COL"=1)

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The estimated cardinality has dropped to 10

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Pagination query using analytic functions

```
select /*+ first_rows(10) */
      *
from   (
       select
          row_number() over (order by sort_order) as rn
          , t.*
       from
          first_k_rows_test t
       where
          filter_col = 1
       order by
          sort_order
      )
where
      rn between 11 and 20
-- and      rownum <= 10 -- Instead of FIRST_ROWS(10) hint
order by
      rn;
```

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Analytic functions allow to write the pagination query in a more expressive way. However they require either the `FIRST_ROWS(n)` hint or adding another `ROWNUM` predicate to activate the optimizer mode - there is no automatic detection similar to `ROWNUM` for the analytic function.

Note that in this case we can safely use the final `ROWNUM` predicate on the same level as the `ORDER BY` since the filter predicate on `RN` already identifies the correct rows

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan slightly different

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	13 (8)
1	SORT ORDER BY		10	13 (8)
* 2	VIEW		10	12 (0)
* 3	WINDOW NOSORT STOPKEY		10	12 (0)
4	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	10	12 (0)
* 5	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX		2 (0)

Predicate Information (identified by operation id):

```
2 - filter("RN">=11 AND "RN"<=20)
3 - filter(ROW_NUMBER() OVER ( ORDER BY "T"."SORT_ORDER")<=20)
5 - access("FILTER_COL"=1)
```

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It's interesting to note that the filter predicates suggest that the optimizer should be able to work out the FIRST\_ROWS(n) scenario automatically. Also note the discrepancy between the filter predicate (ROW\_NUMBER()... <= 20) and the estimates ROWS.

Another point to consider is that the ORDER BY is mandatory to get the FIRST\_ROWS(n) mode to work - although the ROW\_NUMBER() analytic function implies the sort operation already

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Effect of FIRST\_ROWS(n) hint

```
select /*+ first_rows(10) */
      *
from   (
      select
            row_number() over (order by sort_order) as rn
            , t.*
      from
            first_k_rows_test t
      where
            filter_col = 1
      order by
            sort_order
      )
where
      rn between 3991 and 4000
-- and      rownum <= 10 -- Instead of FIRST_ROWS(10) hint
order by
      rn;
```

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The Analytic Function approach has a disadvantage - using the FIRST\_ROWS(n) hint (or the final ROWNUM) confuses the optimizer - see the estimates and actual execution on the next slide

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

Optimizer is "clueless" - actual execution

Id	Operation	Name	Cost	E-Rows	A-Rows	Buffers
1	SORT ORDER BY		13	10	10	3984
* 2	VIEW		12	10	10	3984
* 3	WINDOW NOSORT STOPKEY		12	10	4000	3984
4	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	12	10	4001	3984
* 5	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX	2		4001	15

Predicate Information (identified by operation id):

```
-----  
2 - filter(("RN">=3991 AND "RN"<=4000))  
3 - filter(ROW_NUMBER() OVER ( ORDER BY "T"."SORT_ORDER")<=4000)  
5 - access("FILTER_COL"=1)
```

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The estimated cost clearly doesn't reflect what is happening at runtime

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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## Things look better with ROWNUM

```
select
  *
from (
  select
    rownum as rn
  , v1.*
  from (
    select
      t.*
    from
      first_k_rows_test t
    where
      filter_col = 1
    order by
      sort_order
  ) v1
  where
    rownum <= 4000
  ) v2
where rn >= 3991
order by rn;
```

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Requesting further pages with the ROWNUM approach makes the optimizer aware of the additional work required.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		4000	4226 (1)
1	SORT ORDER BY		4000	4226 (1)
* 2	VIEW		4000	4017 (1)
* 3	COUNT STOPKEY			
4	VIEW		4000	4017 (1)
5	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	4000	4017 (1)
* 6	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX		14 (0)

Predicate Information (identified by operation id):

- 2 - filter("RN">=3991)
- 3 - filter(ROWNUM<=4000)
- 6 - access("FILTER\_COL"=1)

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Here the cardinality estimate is better, but as already shown the optimizer doesn't get the pagination right and therefore believes it has to sort 4000 rows.

This example also makes another point very obvious: The standard Top N / Pagination queries are working well when requesting only the first pages, but the work required increases with the page number requested.

Is there a way to minimize the work and ideally have the same amount of work regardless of the page requested?

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

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Advanced Pagination query - see next page

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The purpose of this re-write is to avoid the costly access to the table for all the rows that will be discarded since they are prior to the requested page. This is achieved here using an inline view that only selects the ROWID from the original pagination query. This can be performed using the index only - a join is then performed to the original table using the ROWID.

Note that Oracle performs a similar operation "under the covers" when using sorted hash clusters - it gets the ROWIDs, sorts them and then gets the maining data from the cluster. This avoids large sort operations, but has the potential overhead of getting single rows by ROWIDs.

This way the performance ought to be almost the same independent from the number of pages requested.

The additional ROWNUM predicate limiting the inner set to 10 rows is necessary here, otherwise we see again that the optimizer doesn't recognize the "pagination" and estimates 4000 rows to join with the original table - eventually leading to a different join mechanism used with increasing number of rows.

```

select
    i.rn
    , t.*
from
    (
        select
            *
        from
            (
                select
                    rownum as rn
                    , v1.*
                from
                    (
                        select
                            t.rowid as row_id
                        from
                            first_k_rows_test t
                        where
                            filter_col = 1
                        order by
                            sort_order
                    ) v1
                where
                    rownum <= 4000
            )
        where
            rn >= 3991
            and rownum <= 10
    ) i
    , first_k_rows_test t
where
    t.rowid = i.row_id
order by
    rn;

```

53

The purpose of this re-write is to avoid the costly access to the table for all the rows that will be discarded since they are prior to the requested page. This is achieved here using an inline view that only selects the ROWID from the original pagination query. This can be performed using the index only - a join is then performed to the original table using the ROWID.

Note that Oracle performs a similar operation "under the covers" when using sorted hash clusters - it gets the ROWIDs, sorts them and then gets the maining data from the cluster. This avoids large sort operations, but has the potential overhead of getting single rows by ROWIDs.

This way the performance ought to be almost the same independent from the number of pages requested.

The additional ROWNUM predicate limiting the inner set to 10 rows is necessary here, otherwise we see again that the optimizer doesn't recognize the "pagination" and estimates 4000 rows to join with the original table - eventually leading to a different join mechanism used with increasing number of rows.

```

select
    i.rn
    , t.*
from
    (
        select
            *
        from
            (
                select
                    rownum as rn
                    , v1.*
                from
                    (
                        select
                            t.rowid as row_id
                        from
                            first_k_rows_test t
                        where
                            filter_col = 1
                        order by
                            sort_order
                    ) v1
                where
                    rownum <= 4000
            )
        where
            rn >= 3991
            and rownum <= 4000
        ) i
    , first_k_rows_test t
where
    t.rowid = i.row_id
order by
    rn;

```

Get the first 4000 rows, ROWID only, therefore index only operation

54

This is achieved here using an inline view that only selects the ROWID from the original pagination query. This can be performed using the index only

Discard all rows (ROWIDs) of the previous pages except for the page requested

```
select
  i.rn
  , t.*
from
  (
    select
      *
    from
      (
        select
          rownum as rn
          , vl.*
        from
          (
            select
              t.rowid as row_id
            from
              first_k_rows_test t
            where
              filter_col = 1
            order by
              sort_order
          ) vl
        where
          rownum <= 4000
      )
    where
      rn >= 3991
    and
      rownum <= 10
  ) i
  , first_k_rows_test t
where
  t.rowid = i.row_id
order by
  rn;
```

55

All rows (effectively ROWIDs only) not required are discarded in the next step

The additional ROWNUM predicate limiting the inner set to 10 rows is necessary here, otherwise we see again that the optimizer doesn't recognize the "pagination" and estimates 4000 rows to join with the original table - eventually leading to a different join mechanism used with increasing number of rows.

Finally pick only the rows from the table that belong to the page requested - minimizing the necessary visits to the table itself

```
select
  i.rn
  , t.*
from
  (
    select
      *
    from
      (
        select
          rownum as rn
          , vl.*
        from
          (
            select
              t.rowid as row_id
            from
              first_k_rows_test t
            where
              filter_col = 1
            order by
              sort_order
          ) vl
        where
          rownum <= 4000
      )
    where
      rn >= 3991
      and rownum <= 10
  ) i
  , first_k_rows_test t
where
  t.rowid = i.row_id
order by
  rn;
```

56

Finally a join is then performed to the original table using the ROWID.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan

Id	Operation	Name	Starts	E-Rows	A-Rows	Buffers
1	SORT ORDER BY		1	10	10	25
2	NESTED LOOPS		1	10	10	25
3	VIEW		1	10	10	15
* 4	COUNT STOPKEY		1		10	15
* 5	VIEW		1	4000	10	15
* 6	COUNT STOPKEY		1		4000	15
7	VIEW		1	4718	4000	15
* 8	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX	1	4718	4000	15
9	TABLE ACCESS BY USER ROWID	FIRST_K_ROWS_TEST	10	1	10	10

Predicate Information (identified by operation id):

- 4 - filter(ROWNUM<=10)
- 5 - filter("RN">=3991)
- 6 - filter(ROWNUM<=4000)
- 8 - access("FILTER\_COL"=1)

57

Compare this to the simple pagination query - the number of buffer gets is much less.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

---

Advanced Pagination query using subquery  
see next page

58

This re-write is based on the assumption that the order expression is unique. In this case we can use a subquery to identify the first row of our page.

This can be done using the existing index as an index-only operation and is therefore very efficient.

The outer query then simply retrieves the n rows of the page again using the same index using the row determined to be the first one as start.

```

select
    3991 + rownum - 1 as rn
, t.*
from (
    select
        t.*
    from
        first_k_rows_test t
    where
        filter_col = 1
    and
        sort_order >=
        (
            select
                sort_order
            from (
                select
                    sort_order
                    , rownum rn
                from (
                    select
                        sort_order
                    from
                        first_k_rows_test t
                    where
                        filter_col = 1
                    order by
                        sort_order
                )
                where rownum <= 3991
            )
            where rn = 3991
        )
    order by sort_order
) t
where rownum <= 10;

```

59

This re-write is based on the assumption that the order expression is unique. In this case we can use a subquery to identify the first row of our page.

This can be done using the existing index as an index-only operation and is therefore very efficient.

The outer query then simply retrieves the n rows of the page again using the same index using the row determined to be the first one as start.

Get the  
SORT\_ORDER for  
the first 3991  
rows, index-only

```
select
  3991 + rownum - 1 as rn
, t.*
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  and
    sort_order >=
      (
        select
          sort_order
        from (
          select
            sort_order
            , rownum rn
          from (
            select
              sort_order
            from
              first_k_rows_test t
            where
              filter_col = 1
            order by
              sort_order
          )
          where rownum <= 3991
        )
      )
  where rn = 3991
)
order by sort_order
) t
where rownum <= 10;
```

60

The innermost query identifies the order expression of all the rows up to the first row of the requested page. This can be done as an index-only operation in our case.

Discard all rows except for the first row of the requested page

```
select
  3991 + rownum - 1 as rn
, t.*
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  and
    sort_order >=
      (
        select
          sort_order
        from (
          select
            sort_order
            , rownum rn
          from (
            select
              sort_order
            from
              first_k_rows_test t
            where
              filter_col = 1
            order by
              sort_order
          )
          where rownum <= 3991
        )
        where rn = 3991
      )
  order by sort_order
) t
where rownum <= 10;
```

61

All rows will be discarded except for the first row of the identified order expressions

Get the actual rows from the table, using the same filter expression again, but starting with the first row identified in the subquery

```
select
  3991 + rownum - 1 as rn
, t.*
from
  (
select
  t.*
from
  first_k_rows_test t
where
  filter_col = 1
and
  sort_order >=
  (
select
  sort_order
from
  (
select
  sort_order
, rownum rn
from
  (
select
  sort_order
from
  first_k_rows_test t
where
  filter_col = 1
order by
  sort_order
)
where rownum <= 3991
)
where rn = 3991
)
order by sort_order
) t
where rownum <= 10;
```

62

The outer query then simply retrieves the n rows of the page again using the same index using the row determined to be the first one as start.

Limit the number of rows accessed to the page size - so we need to visit only as many table rows as the page size

```
select
    3991 + rownum - 1 as rn
, t.*
from
(
    select
        t.*
    from
        first_k_rows_test t
    where
        filter_col = 1
    and
        sort_order >=
        (
            select
                sort_order
            from
                (
                    select
                        sort_order
                        , rownum rn
                    from
                        (
                            select
                                sort_order
                            from
                                first_k_rows_test t
                            where
                                filter_col = 1
                            order by
                                sort_order
                        )
                        where rownum <= 3991
                )
                where rn = 3991
            )
        order by sort_order
    ) t
where rownum <= 10;
```

63

So again the number of rows that need to be obtained from the actual table are limited to the number of rows per page.

# Cost-Based Optimizer (CBO) - Top N - Pagination Queries

## Execution plan

Id	Operation	Name	Starts	E-Rows	A-Rows	Buffers
* 1	COUNT STOPKEY		1		10	27
2	VIEW		1	10	10	27
3	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	1	10	10	27
* 4	<b>INDEX RANGE SCAN</b>	<b>FIRST_K_ROWS_TEST_IDX</b>	1	42	10	17
* 5	VIEW		1	200	1	15
* 6	COUNT STOPKEY		1		3991	15
7	VIEW		1	200	3991	15
* 8	<b>INDEX RANGE SCAN</b>	<b>FIRST_K_ROWS_TEST_IDX</b>	1	200	3991	15

Predicate Information (identified by operation id):

```

1 - filter(ROWNUM<=10)
4 - access("FILTER_COL"=1 AND "SORT_ORDER">= AND "SORT_ORDER" IS NOT NULL)
5 - filter("RN"=3991)
6 - filter(ROWNUM<=3991)
8 - access("FILTER_COL"=1)

```

64

As you can see - although the query looks much more complicated - it is almost as efficient as the one joining by ROWID - the only difference is the second access to the index which adds only a very small overhead (two buffer gets in this example here to retrieve the 10 rows of the page).

One advantage of this query is that we don't need to join by ROWID but can simply use the existing index (twice).

# Everything You Wanted To Know About FIRST\_ROWS\_n

---

A deeper look at  
FIRST\_ROWS(\_n),  
examples, quirks, oddities

# A Deeper Look - FIRST\_ROWS mechanics

---

## Overview of FIRST\_ROWS mode (1)

- Two different scenarios
  - Single table access
  - Joins
- `_sort_elimination_cost_ratio` defines the "breakpoint" between index usage to avoid sort and plans with a sort operation

# A Deeper Look - FIRST\_ROWS mechanics

---

## Overview of FIRST\_ROWS mode (2)

- `_sort_elimination_cost_ratio` default value = 0, see example next slides
- When setting it to a non-default, positive integer then the index to avoid the sort is used if the cost of using it is less than the cost of doing the sort multiplied by the value of the parameter

# A Deeper Look - FIRST\_ROWS mechanics

---

```
_sort_elimination_cost_ratio = 0
```

## Example Setup (1)

```
create table first_k_rows_test
(
  sort_order  varchar2(10) not null,
  filter_col  number        null,
  filler      char(200)
);

insert into first_k_rows_test
select  dbms_random.string('x', 10) as sort_order,
        case when level <= 120 then 1 else 0 end as filter_col,
        'x' as filler
from dual
connect by level <= 100000;

commit;
```

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# A Deeper Look - FIRST\_ROWS mechanics

---

## Example Setup (2)

```
create index first_k_rows_test_idx2 on first_k_rows_test(sort_order);  
  
exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',  
method_opt=>'for all columns size 1 for columns filter_col size 254');
```

# A Deeper Look - FIRST\_ROWS mechanics

---

## ALL\_ROWS Baseline

```
select /*+ all_rows */
      *
from   first_k_rows_test
where  filter_col = 1
order by
      sort_order;
```

# A Deeper Look - FIRST\_ROWS mechanics

## ALL\_ROWS Baseline Plan

```
-----  
| Id | Operation          | Name                | Rows  | Cost (%CPU)|  
-----  
|  0 | SELECT STATEMENT   |                     |  120 |    835  (2)|  
|  1 |   SORT ORDER BY   |                     |    120 |    835  (2)|  
|*  2 |    TABLE ACCESS FULL| FIRST_K_ROWS_TEST |    120 |    834  (2)|  
-----
```

Predicate Information (identified by operation id):

```
-----  
2 - filter("FILTER_COL"=1)
```

# A Deeper Look - FIRST\_ROWS mechanics

---

## FIRST\_ROWS query

```
select /*+ first_rows */
      *
from
      first_k_rows_test
where
      filter_col = 1
order by
      sort_order;
```

# A Deeper Look - FIRST\_ROWS mechanics

## FIRST\_ROWS Plan

```
-----  
| Id | Operation          | Name                | Rows  | Cost (%CPU)|  
-----  
|  0 | SELECT STATEMENT   |                     |  120 |    835  (2)|  
|  1 |   SORT ORDER BY    |                     |    120 |    835  (2)|  
|*  2 |    TABLE ACCESS FULL| FIRST_K_ROWS_TEST  |    120 |    834  (2)|  
-----
```

Predicate Information (identified by operation id):

```
-----  
2 - filter("FILTER_COL"=1)
```

# A Deeper Look - FIRST\_ROWS mechanics

---

## Example Setup (3) - slight modification

```
truncate table first_k_rows_test;

insert into first_k_rows_test
  select dbms_random.string('x', 10) as sort_order,
         case when level <= 121 then 1 else 0 end as filter_col,
         'x' as filler
  from dual
 connect by level <= 100000;

commit;

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',
method_opt=>'for all columns size 1 for columns filter_col size 254');
```

# A Deeper Look - FIRST\_ROWS mechanics

## FIRST\_ROWS Plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		121	100K (1)
* 1	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	121	100K (1)
2	INDEX FULL SCAN	FIRST_K_ROWS_TEST_IDX2	100K	426 (2)

Predicate Information (identified by operation id):

1 - filter("FILTER\_COL"=1)

$120 * 835 = 100,200$

$121 * 835 = 101,035$

`_sort_elimination_cost_ratio = num_rows`

75

It looks like that in recent releases the former behaviour of treating the default setting of `_sort_elimination_cost_ratio = 0` as "infinite" cost for sort operations and therefore avoiding the sort at all costs has changed to setting it to the number of rows estimated to be returned by the query. However this seems only to apply to single-table queries.

# A Deeper Look - FIRST\_ROWS mechanics

---

## Example Setup (4) - Joins

```
truncate table first_k_rows_test;

insert into first_k_rows_test
  select dbms_random.string('x', 10) as sort_order,
         case when level <= 10 then 1 else 0 end as filter_col,
         'x' as filler
  from dual
 connect by level <= 100000;

commit;

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',
method_opt=>'for all columns size 1 for columns filter_col size 254');
```

# A Deeper Look - FIRST\_ROWS mechanics

---

## ALL\_ROWS Baseline

```
select /*+ all_rows */
      *
from
      first_k_rows_test a,
      first_k_rows_test b
where
      a.filter_col = 1
and    a.sort_order = b.sort_order
order by
      a.sort_order;
```

# A Deeper Look - FIRST\_ROWS mechanics

## ALL\_ROWS Baseline Plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	855 (2)
1	SORT ORDER BY		10	855 (2)
2	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	1	2 (0)
3	NESTED LOOPS		10	854 (2)
* 4	TABLE ACCESS FULL	FIRST_K_ROWS_TEST	10	834 (2)
* 5	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX2	1	1 (0)

Predicate Information (identified by operation id):

- 4 - filter("A"."FILTER\_COL"=1)
- 5 - access("A"."SORT\_ORDER"="B"."SORT\_ORDER")

# A Deeper Look - FIRST\_ROWS mechanics

---

## FIRST\_ROWS query

```
select /*+ first_rows */
      *
from
      first_k_rows_test a,
      first_k_rows_test b
where
      a.filter_col = 1
and    a.sort_order = b.sort_order
order by
      a.sort_order;
```

# A Deeper Look - FIRST\_ROWS mechanics

## FIRST\_ROWS Plan

Id	Operation	Name	Rows	Cost (%CPU)
0	SELECT STATEMENT		10	100K (1)
1	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	1	2 (0)
2	NESTED LOOPS		10	100K (1)
* 3	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST	10	100K (1)
4	INDEX FULL SCAN	FIRST_K_ROWS_TEST_IDX2	100K	426 (2)
* 5	INDEX RANGE SCAN	FIRST_K_ROWS_TEST_IDX2	1	1 (0)

Predicate Information (identified by operation id):

3 - filter("A"."FILTER\_COL"=1)  
5 - access("A"."SORT\_ORDER"="B"."SORT\_ORDER")

So we are back at old behaviour with joins

**`_sort_elimination_cost_ratio = 0` means  
infinite**

80

If joins are involved then I couldn't identify any similar "breakpoints" - it looks like the old behaviour of treating 0 as "infinite" is used in this scenario.

# A Deeper Look - FIRST\_ROWS\_n mechanics

## Overview of FIRST\_ROWS\_n mode

- Controlled via `_optimizer_fkr_index_cost_bias`
- Default value = 10
- Range of values: 1..1000
- ```
index_cost <
greatest(
  FTS_cost,
  reduced FTS_cost * _optimizer_fkr_index_cost_bias
)
```
- But... ?

81

The FIRST\_ROWS\_n also has at least one internal parameter that affects the calculation. It seems however not to be limited to the "sort elimination" but obviously controls the "bias" towards index access.

In general above formula seems to apply, which means that an index will be used if its cost is less than the cost of the full table scan multiplied by the parameter.

Since the FIRST\_ROWS\_n mode "knows" the number of rows requested it can actually scale the full table scan cost accordingly - based on the assumption that you don't need to visit all rows from the table to obtain the first n rows.

Of course this assumption is only valid if a SORT is not involved - otherwise you need to obtain all rows identified to perform a sort operation on those rows, or you need to use an index to avoid the sort operation.

There seems to be a sanity check built into the calculation: If the "reduced" FTS cost multiplied by the parameter is less than the cost of the non-scaled FTS cost, then the non-scaled FTS cost is used as lower limit.

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## Test setup (1)

```
create table t2
pctfree 40
pctused 1
as
select
    mod(id, 10) + 1 as col1
    , id as col2
    , id as colx
    , rpad('x', 1000, 'x') as filler
from
    (
        select
            level as id
        from
            dual
        connect by
            level <= 1000
    );
```

82

Here are some examples that do not adhere to the formula provided - so there is obviously something influencing the calculation that I haven't identified yet

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## Test setup (1)

```
update t2 set col1 = 1 where col1 = 2 and rownum <= 1;  
  
commit;  
  
create index t2_idx_filter on t2 (col1);  
  
create index t2_idx_sort on t2 (col1, col2);  
  
exec dbms_stats.gather_table_stats(null, 't2',  
    method_opt=>'for all columns size 1 for columns col1 size  
    254')
```

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The table has 1,000 rows and 100 rows from 1 to 10 in COL1. However I change the distribution slightly - value 1 will be there 101 times, whereas value 2 only 99 times.

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## ALL\_ROWS result

```
select * from t2 where col1 = 1;
```

| Id  | Operation         | Name | Rows | Cost (%CPU) |
|-----|-------------------|------|------|-------------|
| 0   | SELECT STATEMENT  |      | 101  | 69 (0)      |
| * 1 | TABLE ACCESS FULL | T2   | 101  | 69 (0)      |

```
select * from t2 where col1 = 2;
```

| Id  | Operation         | Name | Rows | Cost (%CPU) |
|-----|-------------------|------|------|-------------|
| 0   | SELECT STATEMENT  |      | 99   | 69 (0)      |
| * 1 | TABLE ACCESS FULL | T2   | 99   | 69 (0)      |

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# A Deeper Look - FIRST\_ROWS\_n mechanics

```
_optimizer_fkr_index_cost_bias = 1
```

```
select * from (  
  select * from t2 where coll = 1  
)  
where rownum <= 67;
```

| Id  | Operation                   | Name          | Rows | Cost (%CPU) |
|-----|-----------------------------|---------------|------|-------------|
| 0   | SELECT STATEMENT            |               | 67   | 69 (0)      |
| * 1 | COUNT STOPKEY               |               |      |             |
| 2   | TABLE ACCESS BY INDEX ROWID | T2            | 67   | 69 (0)      |
| * 3 | INDEX RANGE SCAN            | T2_IDX_FILTER |      | 1 (0)       |

```
select * from (  
  select * from t2 where coll = 1  
)  
where rownum <= 68;
```

| Id  | Operation         | Name | Rows | Cost (%CPU) |
|-----|-------------------|------|------|-------------|
| 0   | SELECT STATEMENT  |      | 68   | 48 (0)      |
| * 1 | COUNT STOPKEY     |      |      |             |
| * 2 | TABLE ACCESS FULL | T2   | 68   | 48 (0)      |

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We can clearly see the breakpoint here when the cost exceeds the cost of the full table scan in ALL\_ROWS mode

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

```
_optimizer_fkr_index_cost_bias = 1
```

```
select * from (  
select * from t2 where col1 = 2  
)  
where rownum <= 98;
```

| Id  | Operation                   | Name          | Rows | Cost (%CPU) |
|-----|-----------------------------|---------------|------|-------------|
| 0   | SELECT STATEMENT            |               | 98   | 100 (0)     |
| * 1 | COUNT STOPKEY               |               |      |             |
| 2   | TABLE ACCESS BY INDEX ROWID | T2            | 98   | 100 (0)     |
| * 3 | INDEX RANGE SCAN            | T2_IDX_FILTER |      | 1 (0)       |

86

But it doesn't apply to this case obviously...

# A Deeper Look - FIRST\_ROWS\_n mechanics

```
_optimizer_fkr_index_cost_bias = 1
```

```
select * from (  
select * from t2 where col1 = 2  
order by col2  
)  
where rownum <= 67;
```

| Id  | Operation                   | Name        | Rows | Cost (%CPU) |
|-----|-----------------------------|-------------|------|-------------|
| 0   | SELECT STATEMENT            |             | 67   | 70 (0)      |
| * 1 | COUNT STOPKEY               |             |      |             |
| 2   | VIEW                        |             | 67   | 70 (0)      |
| 3   | TABLE ACCESS BY INDEX ROWID | T2          | 99   | 70 (0)      |
| * 4 | INDEX RANGE SCAN            | T2_IDX_SORT | 67   | 2 (0)       |

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Adding a sort operation gets us back to the expected behaviour

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

```
_optimizer_fkr_index_cost_bias = 1
```

```
select * from (  
select * from t2 where col1 = 2  
order by col2  
)  
where rownum <= 68;
```

| Id  | Operation             | Name | Rows | Cost (%CPU) |
|-----|-----------------------|------|------|-------------|
| 0   | SELECT STATEMENT      |      | 68   | 70 (2)      |
| * 1 | COUNT STOPKEY         |      |      |             |
| 2   | VIEW                  |      | 99   | 70 (2)      |
| * 3 | SORT ORDER BY STOPKEY |      | 99   | 70 (2)      |
| * 4 | TABLE ACCESS FULL     | T2   | 99   | 69 (0)      |

# A Deeper Look - FIRST\_ROWS\_n mechanics

`_optimizer_fkr_index_cost_bias = 1000 ????`

```
select * from (  
select * from t2 where col1 = 2  
order by col2  
)  
where rownum <= 68;
```

| Id  | Operation             | Name | Rows | Cost (%CPU) |
|-----|-----------------------|------|------|-------------|
| 0   | SELECT STATEMENT      |      | 68   | 70 (2)      |
| * 1 | COUNT STOPKEY         |      |      |             |
| 2   | VIEW                  |      | 99   | 70 (2)      |
| * 3 | SORT ORDER BY STOPKEY |      | 99   | 70 (2)      |
| * 4 | TABLE ACCESS FULL     | T2   | 99   | 69 (0)      |

89

But the index usage can't be forced in this case via the parameter

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## Test setup (2)

```
create table first_k_rows_test
(
  sort_order  varchar2(10),
  filter_col   number,
  filler      char(200)
);

insert into first_k_rows_test
select  dbms_random.string('x', 10),
        case
          when level <= 5000
            then 1
          else 0 end,
        'x' as filler
from dual
connect by level <= 100000;

commit;
```

90

Now the previous example suggests that introducing an ORDER BY made the formular work again - potentially due to the fact that there is no "reduced" FTS cost if a sort operation is involved - all rows of an unsorted source need to be processed in order to find the top N rows.

Here is however another example that uses a SORT ORDER BY and always uses the index, regardless of the parameter setting.

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## Test setup (2)

```
create index first_k_rows_test_idx on first_k_rows_test
(filter_col, sort_order);

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test',
method_opt=>'for all columns size 1 for columns filter_col
size 254', estimate_percent=>null);
```

# A Deeper Look - FIRST\_ROWS\_n mechanics

---

## ALL\_ROWS result

```
select
  t.*
from
  first_k_rows_test t
where
  filter_col = 1
order by
  sort_order;
```

| Id  | Operation         | Name              | Rows | Cost (%CPU) |
|-----|-------------------|-------------------|------|-------------|
| 0   | SELECT STATEMENT  |                   | 5000 | 1071 (2)    |
| 1   | SORT ORDER BY     |                   | 5000 | 1071 (2)    |
| * 2 | TABLE ACCESS FULL | FIRST_K_ROWS_TEST | 5000 | 834 (2)     |

# A Deeper Look - FIRST\_ROWS\_n mechanics

```

_optimizer_fkr_index_cost_bias = 1 ???
select
  *
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  order by
    sort_order
)
where
  rownum <= 4999;

```

| Id  | Operation                   | Name                  | Rows | Cost (%CPU) |
|-----|-----------------------------|-----------------------|------|-------------|
| 0   | SELECT STATEMENT            |                       | 4999 | 5024 (1)    |
| * 1 | COUNT STOPKEY               |                       |      |             |
| 2   | VIEW                        |                       | 5000 | 5024 (1)    |
| 3   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST     | 5000 | 5024 (1)    |
| * 4 | INDEX RANGE SCAN            | FIRST_K_ROWS_TEST_IDX | 18   | (0)         |

93

In this example however the index usage can not be prevented regardless of the setting of the parameter.

Further tests indicate that by reading the 10053 optimizer trace debug file it looks like that even with SORT ORDER BYs involved the calculation is based on the "reduced/scaled down" FTS cost in general, which seems to be odd - there is no "reduced" FTS cost if the source is unsorted.

# A Deeper Look - FIRST\_ROWS\_n

## Requesting more rows than estimated in ALL\_ROWS mode

```
select
    t.*
from
    first_k_rows_test t
where
    filter_col = 1;
```

| Id  | Operation         | Name              | Rows | Cost (%CPU) |
|-----|-------------------|-------------------|------|-------------|
| 0   | SELECT STATEMENT  |                   | 5003 | 834 (2)     |
| * 1 | TABLE ACCESS FULL | FIRST_K_ROWS_TEST | 5003 | 834 (2)     |

94

Requesting more rows than estimated in ALL\_ROWS mode effectively disables the FIRST\_ROWS(n) mode - the statement is optimized using ALL\_ROWS mode

# A Deeper Look - FIRST\_ROWS\_n

```
select
  *
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  order by
    sort_order
)
where
  rownum <= 4999;
```

| Id  | Operation                   | Name                  | Rows | Cost (%CPU) |
|-----|-----------------------------|-----------------------|------|-------------|
| 0   | SELECT STATEMENT            |                       | 4999 | 5029 (1)    |
| * 1 | COUNT STOPKEY               |                       |      |             |
| 2   | VIEW                        |                       | 5003 | 5029 (1)    |
| 3   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST     | 5003 | 5029 (1)    |
| * 4 | INDEX RANGE SCAN            | FIRST_K_ROWS_TEST_IDX |      | 21 (0)      |
|     |                             |                       |      | 95          |

Requesting less than 5003 rows will activate FIRST\_ROWS(n) optimization

# A Deeper Look - FIRST\_ROWS\_n

```
select
  *
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  order by
    sort_order
)
where
  rownum <= 5010;
```

| Id  | Operation             | Name              | Rows | Cost (%CPU) |
|-----|-----------------------|-------------------|------|-------------|
| 0   | SELECT STATEMENT      |                   | 5003 | 1071 (2)    |
| * 1 | COUNT STOPKEY         |                   |      |             |
| 2   | VIEW                  |                   | 5003 | 1071 (2)    |
| * 3 | SORT ORDER BY STOPKEY |                   | 5003 | 1071 (2)    |
| * 4 | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST | 5003 | 834 (2)     |

96

Requesting more rows than estimated in ALL\_ROWS mode effectively disables the FIRST\_ROWS(n) mode - the statement is optimized using ALL\_ROWS mode

# A Deeper Look - FIRST\_ROWS\_n

---

```
alter session set tracefile_identifier = 'first_rows_n_greater_than_all_rows';

alter session set events '10053 trace name context forever, level 1';

explain plan for
select
  *
from (
  select
    t.*
  from
    first_k_rows_test t
  where
    filter_col = 1
  order by
    sort_order
)
where
  rownum <= 5010;

alter session set events '10053 trace name context off';
```

97

This can also be seen from the 10053 optimizer trace debug file where no FIRST\_ROWS(n) specific output can be found after the initial evaluation of the estimated number of rows in ALL\_ROWS mode.

# A Deeper Look - FIRST\_ROWS\_n And DML

## DML and FIRST\_ROWS\_n don't mix

```
update
    first_k_rows_test t
set
    filter_col = 1
where
    filter_col = 0
and
    rownum <= 1
returning
    sort_order into :retval;
```

**Special case: Get just the first row(s) for  
update**

98

Oracle optimizes DML always using the ALL\_ROWS mode regardless of any hints / ROWNUM predicates.

In general this is a good idea since DML must process all rows anyway - however there are special cases where you really would like to get just a single row - an example scenario is getting just one code from a pool codes.

# A Deeper Look - FIRST\_ROWS\_n And DML

## Execution plan

| Id  | Operation         | Name              | Starts | E-Rows | A-Rows | Buffers |
|-----|-------------------|-------------------|--------|--------|--------|---------|
| 1   | UPDATE            | FIRST_K_ROWS_TEST | 1      |        | 1      | 312     |
| * 2 | COUNT STOPKEY     |                   | 1      |        | 1      | 307     |
| * 3 | TABLE ACCESS FULL | FIRST_K_ROWS_TEST | 1      | 90415  | 1      | 307     |

Predicate Information (identified by operation id):

- 2 - filter(ROWNUM<=1)
- 3 - filter("FILTER\_COL"=0)

99

The FULL TABLE SCAN here is definitely not the most efficient access to update just a single row. A range scan on the existing index on FILTER\_COL would be a much better idea.

# A Deeper Look - FIRST\_ROWS\_n And DML

---

## Hints required to get it right

```
update      /**+ index_rs(t) */
            first_k_rows_test t
set
            filter_col = 1
where
            filter_col = 0
and         rownum <= 1
returning
            sort_order into :retval;
```

100

Oracle optimizes DML always using the ALL\_ROWS mode regardless of any hints / ROWNUM predicates.

In general this is a good idea since DML must process all rows anyway - however there are special cases where you really would like to get just a single row - an example scenario is getting just one code from a pool codes.

# A Deeper Look - FIRST\_ROWS\_n And DML

## Execution plan

| Id  | Operation        | Name                  | Starts | E-Rows | A-Rows | Buffers |
|-----|------------------|-----------------------|--------|--------|--------|---------|
| 1   | UPDATE           | FIRST_K_ROWS_TEST     | 1      |        | 1      | 7       |
| * 2 | COUNT STOPKEY    |                       | 1      |        | 1      | 2       |
| * 3 | INDEX RANGE SCAN | FIRST_K_ROWS_TEST_IDX | 1      | 90415  | 1      | 2       |

Predicate Information (identified by operation id):

- 2 - filter(ROWNUM<=1)
- 3 - filter("FILTER\_COL"=0)

101

Notice the much less buffer gets

## A Deeper Look - Bitmap Indexes

---

- FIRST\_ROWS(n) doesn't consider Bitmap Indexes for avoiding ORDER BYs
- FIRST\_ROWS does

# A Deeper Look - Bitmap Indexes

---

## Modified test setup

```
drop index first_k_rows_test_idx2;  
  
create bitmap index first_k_rows_test_bitmap on  
first_k_rows_test (sort_order);  
  
alter table first_k_rows_test  
modify sort_order null;
```

103

The test setup gets just slightly modified - the conventional index is replaced with a bitmap index - and additionally since bitmap indexes also index NULL expressions the column is modified to be nullable.

# A Deeper Look - Bitmap Indexes

## FIRST\_ROWS query

```
select /*+ first_rows */  
 *  
from  
   first_k_rows_test t  
order by  
   sort_order;
```

| Id | Operation                     | Name                            | Rows | Cost (%CPU) |
|----|-------------------------------|---------------------------------|------|-------------|
| 0  | SELECT STATEMENT              |                                 | 100K | 5271 (1)    |
| 1  | TABLE ACCESS BY INDEX ROWID   | FIRST_K_ROWS_TEST               | 100K | 5271 (1)    |
| 2  | BITMAP CONVERSION TO ROWIDS   |                                 |      |             |
| 3  | <b>BITMAP INDEX FULL SCAN</b> | <b>FIRST_K_ROWS_TEST_BITMAP</b> |      |             |

# A Deeper Look - Bitmap Indexes

## FIRST\_ROWS(n) query

```
select * from (  
  select *  
  from      first_k_rows_test t  
  order by  
            sort_order  
)  
where rownum <= 10;
```

| Id  | Operation             | Name              | Rows | Cost (%CPU) |
|-----|-----------------------|-------------------|------|-------------|
| 0   | SELECT STATEMENT      |                   | 10   | 5516 (1)    |
| * 1 | COUNT STOPKEY         |                   |      |             |
| 2   | VIEW                  |                   | 100K | 5516 (1)    |
| * 3 | SORT ORDER BY STOPKEY |                   | 100K | 5516 (1)    |
| 4   | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST | 100K | 833 (2)     |

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# A Deeper Look - Bitmap Indexes

## FIRST\_ROWS(n) query - hinted

```
select * from (  
select /*+ index(t, first_k_rows_test_bitmap) */  
*  
from  
    first_k_rows_test t  
order by  
    sort_order  
)  
where rownum <= 10;
```

| Id  | Operation                     | Name                            | Rows | Cost (%CPU) |
|-----|-------------------------------|---------------------------------|------|-------------|
| 0   | SELECT STATEMENT              |                                 | 10   | 7 (0)       |
| * 1 | COUNT STOPKEY                 |                                 |      |             |
| 2   | VIEW                          |                                 | 10   | 7 (0)       |
| 3   | TABLE ACCESS BY INDEX ROWID   | FIRST_K_ROWS_TEST               | 10   | 7 (0)       |
| 4   | BITMAP CONVERSION TO ROWIDS   |                                 |      |             |
| 5   | <b>BITMAP INDEX FULL SCAN</b> | <b>FIRST_K_ROWS_TEST_BITMAP</b> |      |             |

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# A Deeper Look - Bitmap Indexes

FIRST\_ROWS(n) query - hinted

## Runtime Statistics

| Id  | Operation                   | Name                     | Starts | E-Rows | A-Rows | Buffers |
|-----|-----------------------------|--------------------------|--------|--------|--------|---------|
| * 1 | COUNT STOPKEY               |                          | 1      |        | 10     | 12      |
| 2   | VIEW                        |                          | 1      | 10     | 10     | 12      |
| 3   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST        | 1      | 10     | 10     | 12      |
| 4   | BITMAP CONVERSION TO ROWIDS |                          | 1      |        | 10     | 2       |
| 5   | BITMAP INDEX FULL SCAN      | FIRST_K_ROWS_TEST_BITMAP | 1      |        | 10     | 2       |

# A Deeper Look - Bitmap Indexes

## FIRST\_ROWS(n) query - filter added

```
select * from (  
  select  
    *  
  from  
    first_k_rows_test t  
  where  
    sort_order >= '0'  
  order by  
    sort_order  
)  
where rownum <= 10;
```

| Id  | Operation                   | Name                     | Rows | Cost (%CPU) |
|-----|-----------------------------|--------------------------|------|-------------|
| 0   | SELECT STATEMENT            |                          | 10   | 7 (0)       |
| * 1 | COUNT STOPKEY               |                          |      |             |
| 2   | VIEW                        |                          | 10   | 7 (0)       |
| 3   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST        | 10   | 7 (0)       |
| 4   | BITMAP CONVERSION TO ROWIDS |                          |      |             |
| * 5 | BITMAP INDEX RANGE SCAN     | FIRST_K_ROWS_TEST_BITMAP |      |             |

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## A Deeper Look - Potential parsing overhead

---

- In more recent releases (in particular in 11g and later) there is a measurable parsing overhead when using FIRST\_ROWS(n) mode
- The following block simply hard parses a query 1,000 times

# A Deeper Look - Potential parsing overhead

---

```
declare
  c integer;
begin
  c := dbms_sql.open_cursor;
  for i in 1..1000 loop
    dbms_sql.parse(c, '
select  /* PARSE_' || to_char(i, 'TM') || ' */
      *
from
      first_k_rows_test a
where a.filter_col > (
select
      avg(b.filter_col)
from
      first_k_rows_test b
where
      b.sort_order = a.sort_order
)
      ', dbms_sql.NATIVE);
    end loop;
  dbms_sql.close_cursor(c);
end;
/
```

110

It is a rather simply query with a single subquery

## A Deeper Look - Potential parsing overhead

---

- In ALL\_ROWS mode this takes on my test system approx. **5.5** seconds (11.1.0.7)
- Let's run the same in FIRST\_ROWS(1) mode

# A Deeper Look - Potential parsing overhead

---

```
declare
  c integer;
begin
  c := dbms_sql.open_cursor;
  for i in 1..1000 loop
    dbms_sql.parse(c, '
select   /*+ FIRST_ROWS(1) */ /* PARSE_' || to_char(i, 'TM') || ' */
      *
from
      first_k_rows_test a
where a.filter_col > (
select
      avg(b.filter_col)
from
      first_k_rows_test b
where
      b.sort_order = a.sort_order
)
      ', dbms_sql.NATIVE);
    end loop;
  dbms_sql.close_cursor(c);
end;
/
```

112

## A Deeper Look - Potential parsing overhead

---

- In FIRST\_ROWS(1) mode this takes on my test system approx. **10** seconds (11.1.0.7)
- With well behaving applications not relevant, but there is some overhead that might lead to problems with applications that do a lot of hard parsing

# A Deeper Look - Things Can Go Wrong (1)

---

## Test setup - function-based index

```
create table test_sort(name,object_name)
as
select 'select',object_name from
dba_objects;

create index i_test_sort_1 on
test_sort(substr(name,0,3),object_name);

exec dbms_stats.gather_table_stats(null,'TEST_SORT');
```

114

Oracle 10.2.0.4 has problems when using function-based indexes and first\_rows(n) mode. This might be bug 3940803 which was filed for 9.2.0.5, but there are many bugs in this context, so it could also be a regression of some of these bugs.

# A Deeper Look - Things Can Go Wrong (1)

---

## Simple Top-N query using expression

```
select
    *
from (
    select
        *
    from
        test_sort
    where
        substr(name,0,3) = 'sel'
    order by
        object_name
)
where
    rownum<=10;
```

115

# A Deeper Look - Things Can Go Wrong (1)

## Execution plan 10.2.0.4

```
-----  
| Id | Operation                | Name      | Rows  | Cost (%CPU)|  
-----  
0	SELECT STATEMENT		10	529   (4)
*  1	COUNT STOPKEY			
2	VIEW		50744	529   (4)
*  3	SORT ORDER BY STOPKEY		50744	529   (4)
*  4	TABLE ACCESS FULL	TEST_SORT	50744	75   (6)
-----
```

Predicate Information (identified by operation id):

```
-----  
1 - filter(ROWNUM<=10)  
3 - filter(ROWNUM<=10)  
4 - filter(SUBSTR("NAME",0,3)='sel')
```

116

The function-based index doesn't get used in 10.2.0.4

# A Deeper Look - Things Can Go Wrong (1)

## Execution plan 11.1.0.7

| Id  | Operation                   | Name                 | Rows | Cost (%CPU)  |
|-----|-----------------------------|----------------------|------|--------------|
| 0   | SELECT STATEMENT            |                      | 10   | 8 (0)        |
| * 1 | COUNT STOPKEY               |                      |      |              |
| 2   | VIEW                        |                      | 10   | 8 (0)        |
| 3   | TABLE ACCESS BY INDEX ROWID | TEST_SORT            | 10   | 8 (0)        |
| * 4 | <b>INDEX RANGE SCAN</b>     | <b>I_TEST_SORT_1</b> |      | <b>3</b> (0) |

Predicate Information (identified by operation id):

- 1 - filter(ROWNUM<=10)
- 4 - **access(SUBSTR("NAME",0,3)='sel')**

117

It gets used in 11.1.0.7 however

# A Deeper Look - Things Can Go Wrong (1)

---

## Workaround in 10.2.0.4

```
select
      *
from    (
        select
              *
        from    test_sort
        where   substr(name,0,3) = 'sel'
        order by substr(name,0,3), object_name
        )
where   rownum<=10;
```

118

Adding the expression to the ORDER BY enables the index usage in 10.2.0.4

# A Deeper Look - Things Can Go Wrong (1)

## Workaround execution plan 10.2.0.4

| Id  | Operation                   | Name                 | Rows      | Cost (%CPU)  |
|-----|-----------------------------|----------------------|-----------|--------------|
| 0   | SELECT STATEMENT            |                      | 10        | 8 (0)        |
| * 1 | COUNT STOPKEY               |                      |           |              |
| 2   | VIEW                        |                      | 10        | 8 (0)        |
| 3   | TABLE ACCESS BY INDEX ROWID | TEST_SORT            | 50744     | 8 (0)        |
| * 4 | <b>INDEX RANGE SCAN</b>     | <b>I_TEST_SORT_1</b> | <b>10</b> | <b>3</b> (0) |

Predicate Information (identified by operation id):

1 - filter(ROWNUM<11)  
4 - **access(SUBSTR("NAME",0,3)='sel')**

# A Deeper Look - Things Can Go Wrong (2)

## Test setup - Join with group by and filter

```
create table first_k_rows_test1
as
select
    id
    , mod(id, 100) as id_join
    , mod(id, 100) as filter_col
    , rpad('x', 100) as filler
from
    (
    select
        level as id
    from
        dual
    connect by
        level <= 100000
    );
```

<http://forums.oracle.com/forums/thread.jspa?threadID=934895>

120

This example for an issue with GROUP BY is based on a test case posted on the OTN forums, however the original setup was much more complicated and therefore I've tried to simplify as much as possible.

We have a simple table with 100,000 rows

# A Deeper Look - Things Can Go Wrong (2)

---

## Test setup - Join with group by and filter

```
create table first_k_rows_test2
as
select
    id_join
    , upper(dbms_random.string('a',15)) as group_col
from
    (
        select distinct
            id_join
        from
            first_k_rows_test1
    );

create index first_k_rows_test2_idx on first_k_rows_test2 (group_col);

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test1')

exec dbms_stats.gather_table_stats(null, 'first_k_rows_test2')
```

121

The second table is only a lookup table - the crucial point here is that there is an index on the grouping column.

# A Deeper Look - Things Can Go Wrong (2)

---

## The query

```
select
  group_col
from
  (
    select
      t2.group_col
      , count(*) as cnt
    from
      first_k_rows_test1 t1
      , first_k_rows_test2 t2
    where
      t1.id_join = t2.id_join
      and t2.group_col != 'BLA'
      and t1.filter_col between 1 and 2
    group by
      t2.group_col
    order by
      cnt desc
  )
where
  rownum <= 1;
```

122

Having the index and a filter predicate on the grouping column encourages the optimizer to evaluate an execution plan that avoids the grouping operation.

# A Deeper Look - Things Can Go Wrong (2)

## Execution plan 10.2.0.4 + 11.1.0.7

| Id  | Operation                   | Name                   | Rows | Cost (%CPU) |
|-----|-----------------------------|------------------------|------|-------------|
| 0   | SELECT STATEMENT            |                        | 1    | 35 (3)      |
| * 1 | COUNT STOPKEY               |                        |      |             |
| 2   | VIEW                        |                        | 2    | 35 (3)      |
| * 3 | SORT ORDER BY STOPKEY       |                        | 2    | 35 (3)      |
| 4   | SORT GROUP BY NOSORT        |                        | 2    | 35 (3)      |
| 5   | NESTED LOOPS                |                        | 2    | 35 (3)      |
| 6   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST2     | 99   | 2 (0)       |
| * 7 | INDEX FULL SCAN             | FIRST_K_ROWS_TEST2_IDX | 1    | 1 (0)       |
| * 8 | TABLE ACCESS FULL           | FIRST_K_ROWS_TEST1     | 2    | 33 (4)      |

Predicate Information (identified by operation id):

```
-----  
1 - filter(ROWNUM<=1)  
3 - filter(ROWNUM<=1)  
7 - filter("T2"."GROUP_COL"<>'BLA')  
8 - filter("T1"."FILTER_COL"<=2 AND "T1"."FILTER_COL">=1 AND  
"T1"."ID_JOIN"="T2"."ID_JOIN")
```

123

The problem seems to be that the ROWNUM predicate is pushed far down and therefore the plan is based on the assumption that only a single row needs to read from the small table - however the nested loop will have to run through all iterations to perform the group by operation.

# A Deeper Look - Things Can Go Wrong (2)

## Runtime statistics

| Id  | Operation                   | Name                   | Starts | E-Rows | A-Rows | Buffers |
|-----|-----------------------------|------------------------|--------|--------|--------|---------|
| * 1 | COUNT STOPKEY               |                        | 1      |        | 1      | 164K    |
| 2   | VIEW                        |                        | 1      | 2      | 1      | 164K    |
| * 3 | SORT ORDER BY STOPKEY       |                        | 1      | 2      | 1      | 164K    |
| 4   | SORT GROUP BY NOSORT        |                        | 1      | 2      | 2      | 164K    |
| 5   | NESTED LOOPS                |                        | 1      | 2      | 2000   | 164K    |
| 6   | TABLE ACCESS BY INDEX ROWID | FIRST_K_ROWS_TEST2     | 1      | 99     | 100    | 2       |
| * 7 | INDEX FULL SCAN             | FIRST_K_ROWS_TEST2_IDX | 1      | 1      | 100    | 1       |
| * 8 | TABLE ACCESS FULL           | FIRST_K_ROWS_TEST1     | 100    | 2      | 2000   | 164K    |

Predicate Information (identified by operation id):

```
1 - filter(ROWNUM<=1)
3 - filter(ROWNUM<=1)
7 - filter("T2"."GROUP_COL" <> 'BLA')
8 - filter(("T1"."FILTER_COL" <= 2 AND "T1"."FILTER_COL" >= 1 AND
"T1"."ID_JOIN" = "T2"."ID_JOIN"))
```

124

It is obviously not a good idea to run a full table scan 100 times in a nested loop

## A Deeper Look - Things Can Go Wrong (2)

---

Potential workarounds:

- Use ALL\_ROWS hint
- Use pre-10.2.0.3 optimizer features
- Use "\_first\_k\_rows\_dynamic\_proration" parameter

```
alter session set "_first_k_rows_dynamic_proration" = false;
```

- Obviously functionality changes from release to release, since 10.2.0.3 enabled by default (true)

125

The issue is fixed in 11.2.0.1. Switching to pre-10.2.0.3 optimizer settings simply sets the "\_first\_k\_rows\_dynamic\_proration" parameter to false.

The "\_first\_k\_rows\_dynamic\_proration" seems to enable more sophisticated code paths regarding the FIRST\_ROWS(n) mode - as always new features not always result in better execution plans.

# A Deeper Look - Things Can Go Wrong (2)

## Workaround execution plan

| Id  | Operation             | Name               | Starts | E-Rows | A-Rows | Buffers |
|-----|-----------------------|--------------------|--------|--------|--------|---------|
| * 1 | COUNT STOPKEY         |                    | 1      |        | 1      | 1644    |
| 2   | VIEW                  |                    | 1      | 99     | 1      | 1644    |
| * 3 | SORT ORDER BY STOPKEY |                    | 1      | 99     | 1      | 1644    |
| 4   | HASH GROUP BY         |                    | 1      | 99     | 2      | 1644    |
| * 5 | HASH JOIN             |                    | 1      | 2931   | 2000   | 1644    |
| * 6 | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST2 | 1      | 99     | 100    | 3       |
| * 7 | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST1 | 1      | 2990   | 2000   | 1641    |

Predicate Information (identified by operation id):

- 1 - filter(ROWNUM<=1)
- 3 - filter(ROWNUM<=1)
- 5 - access("T1"."ID\_JOIN"="T2"."ID\_JOIN")
- 6 - filter("T2"."GROUP\_COL"<>'BLA')
- 7 - filter(("T1"."FILTER\_COL"<=2 AND "T1"."FILTER\_COL">=1))

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Notice the much less buffer gets

## A Deeper Look - Things Can Go Wrong (2)

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- But using the "\_first\_k\_rows\_dynamic\_proration" not always helps. Slightly changing the query - for example removing the filter predicate - will result in the bad execution plan, even with the parameter set to "false"
- Let's see how the plan looks like when forcing the nested loop operation and the parameter set to false

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# A Deeper Look - Things Can Go Wrong (2)

## Execution plan with workaround

| Id  | Operation             | Name               | Rows | Cost (%CPU) |
|-----|-----------------------|--------------------|------|-------------|
| 0   | SELECT STATEMENT      |                    | 1    | 458 (3)     |
| * 1 | COUNT STOPKEY         |                    |      |             |
| 2   | VIEW                  |                    | 99   | 458 (3)     |
| * 3 | SORT ORDER BY STOPKEY |                    | 99   | 458 (3)     |
| 4   | HASH GROUP BY         |                    | 99   | 458 (3)     |
| * 5 | HASH JOIN             |                    | 2931 | 456 (2)     |
| * 6 | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST2 | 99   | 2 (0)       |
| * 7 | TABLE ACCESS FULL     | FIRST_K_ROWS_TEST1 | 2990 | 454 (2)     |

Predicate Information (identified by operation id):

```
1 - filter(ROWNUM<=1)
3 - filter(ROWNUM<=1)
5 - access("T1"."ID_JOIN"="T2"."ID_JOIN")
6 - filter("T2"."GROUP_COL"<>'BLA')
7 - filter("T1"."FILTER_COL"<=2 AND "T1"."FILTER_COL">=1)
```

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# A Deeper Look - Things Can Go Wrong (2)

## Execution plan with forced NL

```
-----  
| Id | Operation                               | Name                               | Rows | Cost (%CPU)|  
-----  
0	SELECT STATEMENT		1	463 (3)
* 1	COUNT STOPKEY			
2	VIEW		23	463 (3)
* 3	SORT ORDER BY STOPKEY		23	463 (3)
4	SORT GROUP BY NOSORT		23	463 (3)
5	NESTED LOOPS		23	463 (3)
6	TABLE ACCESS BY INDEX ROWID	FIRST_K_ROWS_TEST2	99	2 (0)
* 7	INDEX FULL SCAN	FIRST_K_ROWS_TEST2_IDX	2	1 (0)
* 8	TABLE ACCESS FULL	FIRST_K_ROWS_TEST1	15	231 (3)
-----
```

Predicate Information (identified by operation id):

```
-----  
1 - filter(ROWNUM<=1)  
3 - filter(ROWNUM<=1)  
7 - filter("T2"."GROUP_COL" <> 'BLA')  
8 - filter("T1"."FILTER_COL" <= 2 AND "T1"."FILTER_COL" >= 1 AND  
"T1"."ID_JOIN" = "T2"."ID_JOIN")
```

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So the plan just got discarded because it is slightly more expensive - a small change to the query will result in that execution plan potentially becoming cheaper.

# Everything You Wanted To Know About FIRST\_ROWS\_n

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## Summary (1)

- Important to understand when to use FIRST\_ROWS(\_n) optimization
- FIRST\_ROWS is deprecated since 9i and might produce suboptimal execution plans, in particular when joins are involved

# Everything You Wanted To Know About FIRST\_ROWS\_n

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## Summary (2)

- Advanced techniques for Pagination queries when frequently accessing later pages
- If there are issues with executions plans, one can try to set the "`_first_k_rows_dynamic_proration`" parameter to false in 10.2.0.4 and later

# Everything You Wanted To Know About FIRST\_ROWS\_n

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Questions

&

Answers

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# Everything You Wanted To Know About FIRST\_ROWS\_n

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Thank you!